



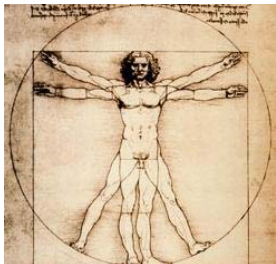
# Hardware Optimization of error control decoders

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**International Workshop on Channel Coding  
and Cooperation in Wireless  
Communications**





**This work has been done in the framework  
of DaVinci project (FP7).**

**With the contributions of :**

**Laura Conde-Canencia**

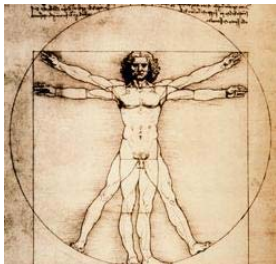
**Ashwani Singh**

**Ali Al-Ghouwayel**

**Pierre Bomel**

**Yang-Yang Tang**

**Cedric Marchand**



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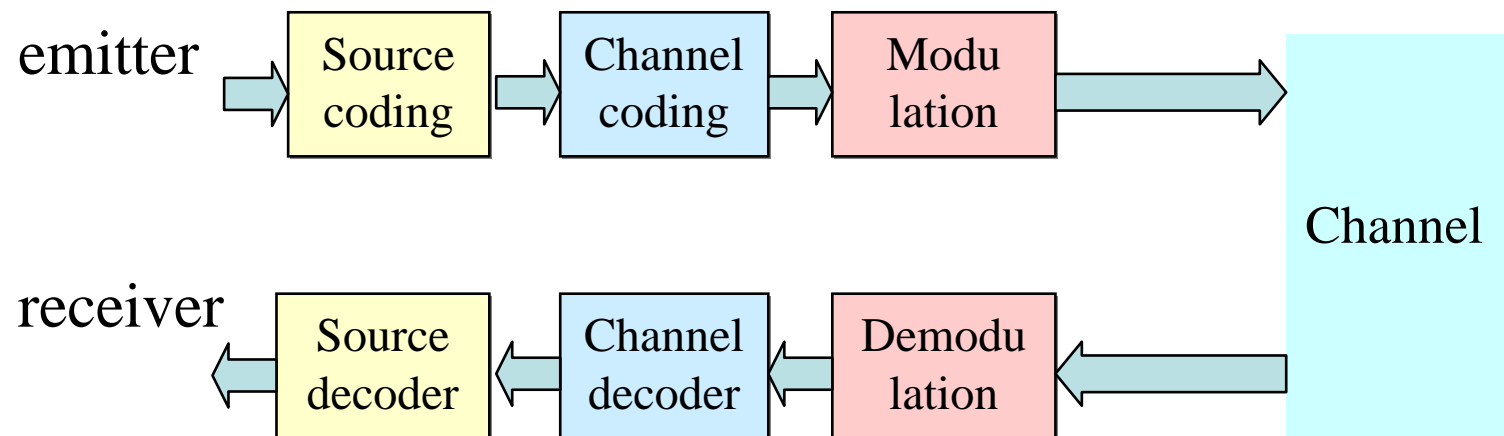
# OUTLINE

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- **MOTIVATION**
- **REDUCED MONTE CARLO SIMULATION**
- **ENTROPY INSPIRED DISTANCE**
- **HARDWARE EMULATION**
- **CONCLUSION**

# Motivation

Design of a communication system...

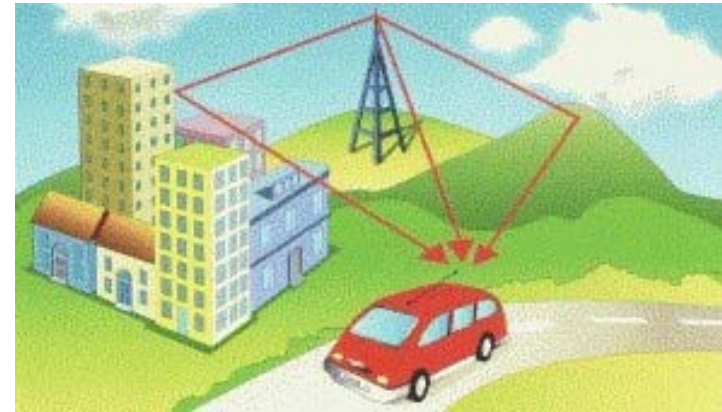


...find the best complexity-performance trade-off

## Motivation

Performance:

- **BER**
- Jammer rejection
- time of synchronization...
- ...



Complexity:

- **area, power dissipation**
- time to market
- ...

**algorithm**  
**ADC resolution,**  
**sampling frequency,**  
**fixed precision**

A very complex problem...

# Monte-carlo simulation

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- Formal expression of the BER: refer to Proakis
- In practice, estimation of the BER using Monte-carlo simulation
  - Software model of emitter, channel, receiver
  - Emulation of the transmission of  $N$  bits
  - Estimation of the BER as  $Nb\_errors/N$

*VERY EFFICIENT...*

*...but*

*TIME CONSUMMING: BER of  $10^{-6}$  (+-3%) requires  $10^9$  bits.*

# Two separate problems

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- 1) The code design problem.
  - Theoretical bound
  - Weight estimation using impulse method and its derivative
  - “Smart” MC simulation (semi-random sample)
  - Estimation of the APP distribution using a reduced set of samples => deduce BER
- 2) Evaluate performance of “hardware decoder” vs “ideal decoder”

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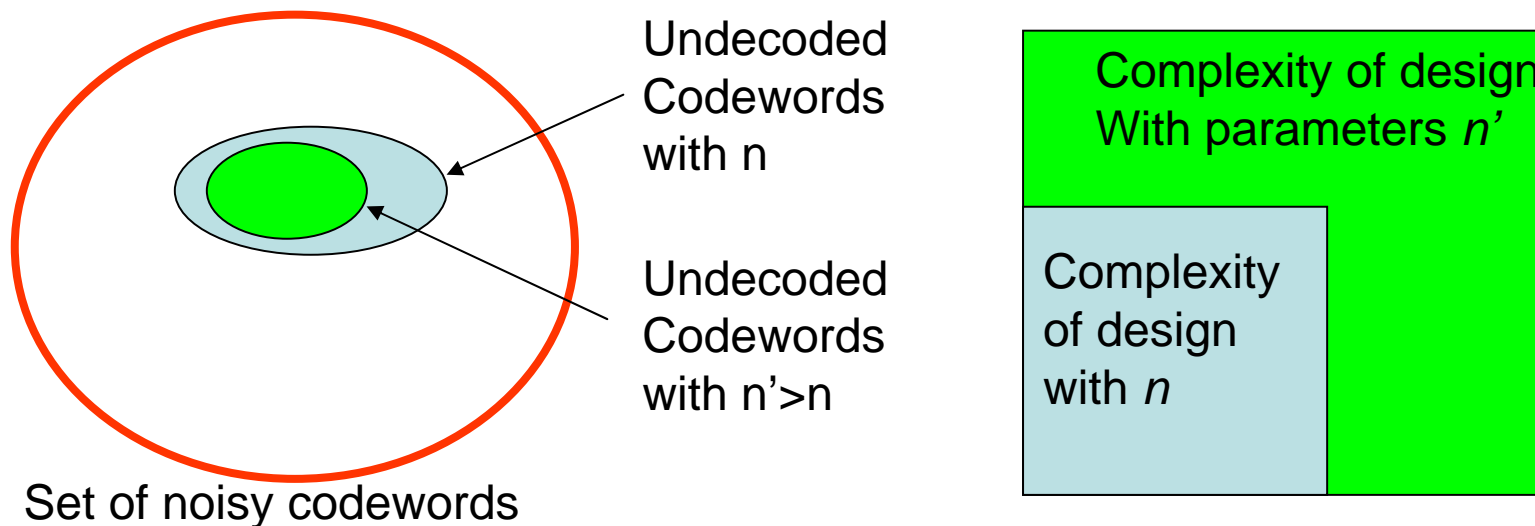
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# REDUCED MONTE CARLO SIMULATION

- Let us assume that we want to find a best trade-off of a algorithm parameter  $n$ . assuming that:
  - 1) Decoder complexity is an increasing function of  $n$ .
  - 2) Performance is a strictly increasing function, i.e. if a noisy codeword is decoded with parameter  $n$ , than the codeword is decoded with parameter  $n' > n$ .



**In that case, RMCS can be efficient**

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# **REDUCED MONTE CARLO SIMULATION**

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- **Step 1: Run Monte-Carlo simulation with the “worst-case” configuration. Store the undecoded codeword in set  $S$  ;**
- **Step 2: Re-run decoder with better parameters only in subset  $S$  to evaluate performances ;**
- **Step 3: perform a classical MCS in order to verify the validity of hypotheses 1 for some configurations.**

# RESULT ON DAVINCI CODE (LDPC GF(64))

10<sup>7</sup> frames

(BER 5.10<sup>-3</sup>)

	<i>it</i> = 8	<i>it</i> = 10	<i>it</i> = 12	<i>it</i> = 14	<i>it</i> = 16
<i>n</i> = 8	49693	13535	5404	2993	2105
<i>n</i> = 10	12340	2569	815	414	291
<i>n</i> = 12	4763	893	270	109	62
<i>n</i> = 14	2419	467	163	72	42
<i>n</i> = 16	1349	223	<del>56</del> 71	22	<del>13</del> 22
<i>n</i> = 18	835	151	25	10	4
<i>n</i> = 20	592	116	17	7	6
<i>n</i> = 22	455	103	14	4	3
<i>n</i> = 24	373	86	<del>14</del> 21	5	<del>4</del> 8
<i>n</i> = 26	321	91	16	4	3
<i>n</i> = 28	262	54	9	5	3
<i>n</i> = 30	235	52	8	4	3
<i>n</i> = 32	228	52	<del>7</del> 19	4	<del>2</del> 8

Save 3 months of simulation...

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# **ENTROPY INSPIRED DISTANCE**

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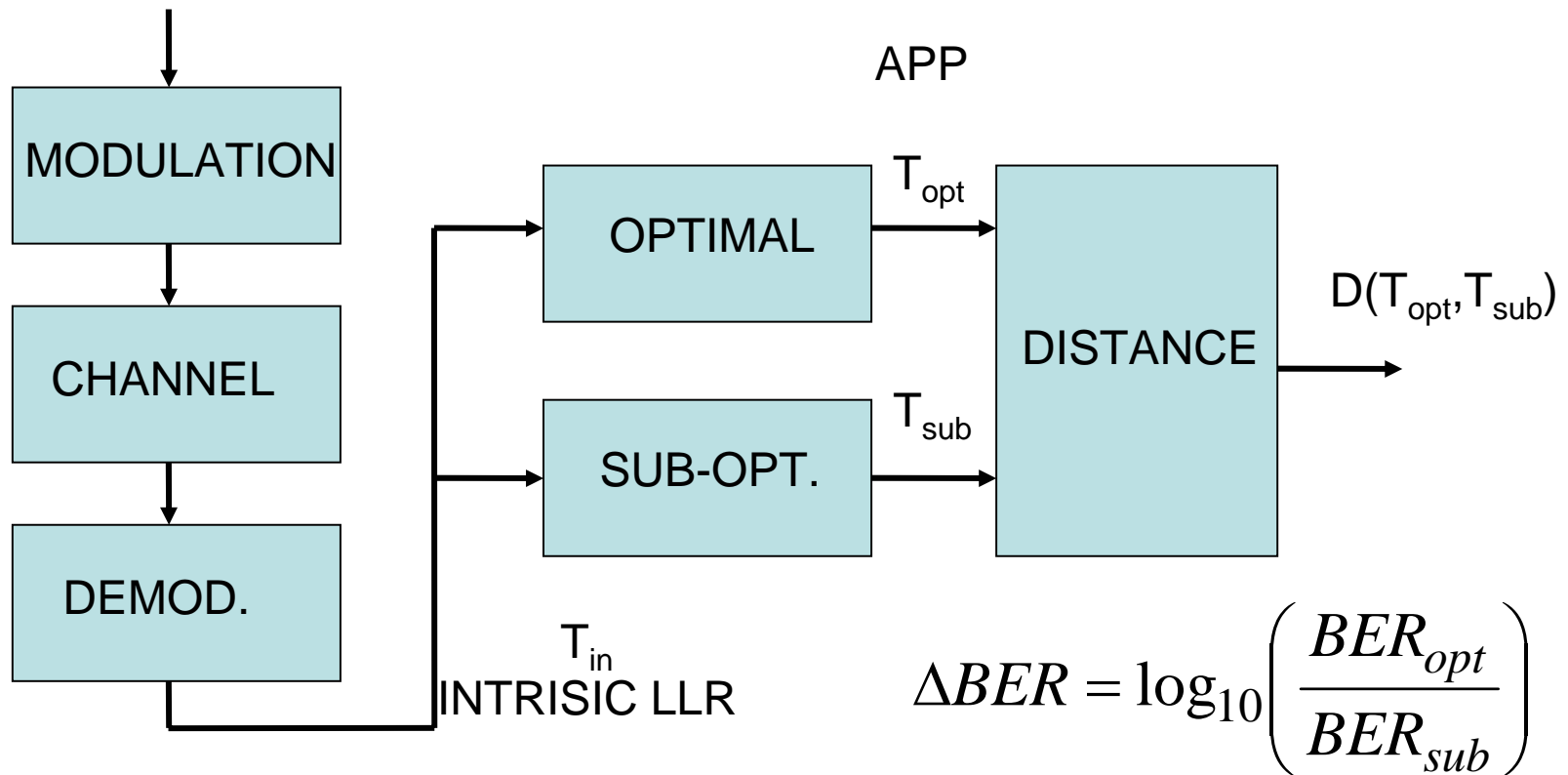
- **CONSTAT : FROM APP TO BER, A HUGE QUANTITY OF INFORMATION IS SUPPRESSED.**

**QUESTION 1 : CAN WE EXPLOIT THIS INFORMATION ?**

**QUESTION 2 : HOW ?**

# DISTANCE BETWEEN OPTIMAL/SUB-OPT.

SET OF  $M$  CODEWORDS OF SIZE  $P$  ( $N=M \times P$   $Q$ -ary symbols)



**FIND  $D(T_{opt}, T_{sub})$  RELATED TO  $\Delta BER$**

## QUESTION : WHAT DISTANCE TO CHOOSE...

- LET  $(p_k^0, p_k^1, \dots, p_k^{Q-1})$  AND  $(\tilde{p}_k^0, \tilde{p}_k^1, \dots, \tilde{p}_k^{Q-1})$  BE THE APP OF THE  $k^{\text{th}}$  DECODED SYMBOL FOR OPTIMAL AND SUB-OPTIMAL DECODERS.

$$\sum_{i=0}^{Q-1} p_k^i = 1$$

$$\sum_{i=0}^{Q-1} \tilde{p}_k^i = 1$$

- FIRST ATTENT : L1-NORM

$$L_1 = \sum_{k=0}^{N-1} \left( \sum_{i=0}^{Q-1} |p_k^i - \tilde{p}_k^i| \right)$$

EXPERIMENTAL RESULT: NO CLEAR CORRELATION BETWEEN  
L1-log(Tsub, Topt) AND  $\Delta$ BER.

PROPOSED EXPLANATION:  $p=10^{-6}$  AND  $p=10^{-12}$  ARE REALLY DIFFERENT  
FROM A DECODING PERSPECTIVE

## QUESTION : WHAT DISTANCE TO CHOOSE...

- LET  $(p_k^0, p_k^1, \dots, p_k^{Q-1})$  AND  $(\tilde{p}_k^0, \tilde{p}_k^1, \dots, \tilde{p}_k^{Q-1})$  BE THE APP OF THE  $k^{\text{th}}$  DECODED SYMBOL FOR OPTIMAL AND SUB-OPTIMAL DECODERS.

$$\sum_{i=0}^{Q-1} p_k^i = 1$$

$$\sum_{i=0}^{Q-1} \tilde{p}_k^i = 1$$

- SECOND ATTENT : L1-log MORM  $L_1 = \sum_{k=0}^{N-1} \left( \sum_{i=0}^{Q-1} -\log |p_k^i - \tilde{p}_k^i| \right)$

EXPERIMENTAL RESULT: NO CLEAR CORRELATION BETWEEN  
L1-log(Tsub, Topt) AND  $\Delta$ BER.

PROPOSED EXPLANATION:  $p=0.7$  AND  $p=0.9$  ARE REALLY DIFFERENT  
FROM A DECODING PERSPECTIVE AND  $-\log(0.2) \ll -\log(10^{-5})$

## QUESTION : WHAT DISTANCE TO CHOOSE...

- LET  $(p_k^0, p_k^1, \dots, p_k^{Q-1})$  AND  $(\tilde{p}_k^0, \tilde{p}_k^1, \dots, \tilde{p}_k^{Q-1})$  BE THE APP OF THE  $k^{\text{th}}$  DECODED SYMBOL FOR OPTIMAL AND SUB-OPTIMAL DECODERS.

$$\sum_{i=0}^{Q-1} p_k^i = 1$$

$$\sum_{i=0}^{Q-1} \tilde{p}_k^i = 1$$

- THIRD ATTENT : KULLBACK-LEIBLER DIVERGENCE

$$KL = \sum_{k=0}^{N-1} \left( \sum_{i=0}^{Q-1} -p_k^i \log(p_k^i / \tilde{p}_k^i) \right)$$

EXPERIMENTAL RESULT: NO CLEAR CORRELATION BETWEEN  
KL(Tsub, Topt) AND  $\Delta$ BER.

PROPOSED EXPLANATION: *NO EXPLANATION...*

# ENTROPY INSPIRED DISTANCE (EID)

- **ENTROPY OF A DISTRIBUTION :**  $H(X) = \sum_{i=0}^{Q-1} -p_k^i \log(p_k^i)$
- **PROPERTIES :**  $H(X) \geq 0$ ,  $H(X)=\log(Q)$  for  $p^i=1/Q$ ,  $i=0..Q-1$ .

$$EID = \sum_{k=0}^{N-1} \left( \sum_{i=0}^{Q-1} - \left| p_k^i - \tilde{p}_k^i \right| \log \left( \left| p_k^i - \tilde{p}_k^i \right| \right) \right)$$

EXPERIMENTAL RESULT: USEFUL TOOLS TO PREDICT THE  $\Delta$ BER !!!

# PROPERTIES OF EID

- IF  $\text{Arg max}(p^i) = \text{Arg max}(\tilde{p}^i)$

THEN EID IS A DISTANCE, I.E.

$$\text{EID}(X, X') = 0 \Rightarrow X = X' \quad (\text{note that } \text{EID}([0, 0, 0, 1], [1, 0, 0, 0]) = 0)$$

$$\text{EID}(X, X') = \text{EID}(X', X)$$

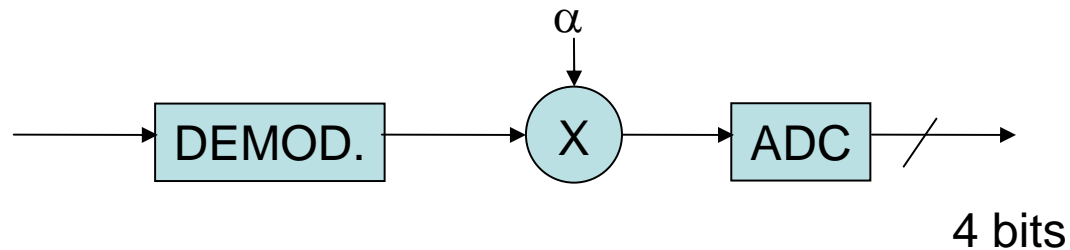
$$\text{EID}(X, Y) + \text{EID}(Y, Z) \geq \text{EID}(X, Z)$$

AT THE MOMENT, PROVEN FOR  $Q=2$  UP TO 12.

...STILL NEED A PROOF FOR  $Q > 12$

# SIMULATION RESULT : CASE 1

- DUO BINARY TURBO-CODE
- SEARCH FOR THE OPTIMAL QUANTIZATION FACTOR

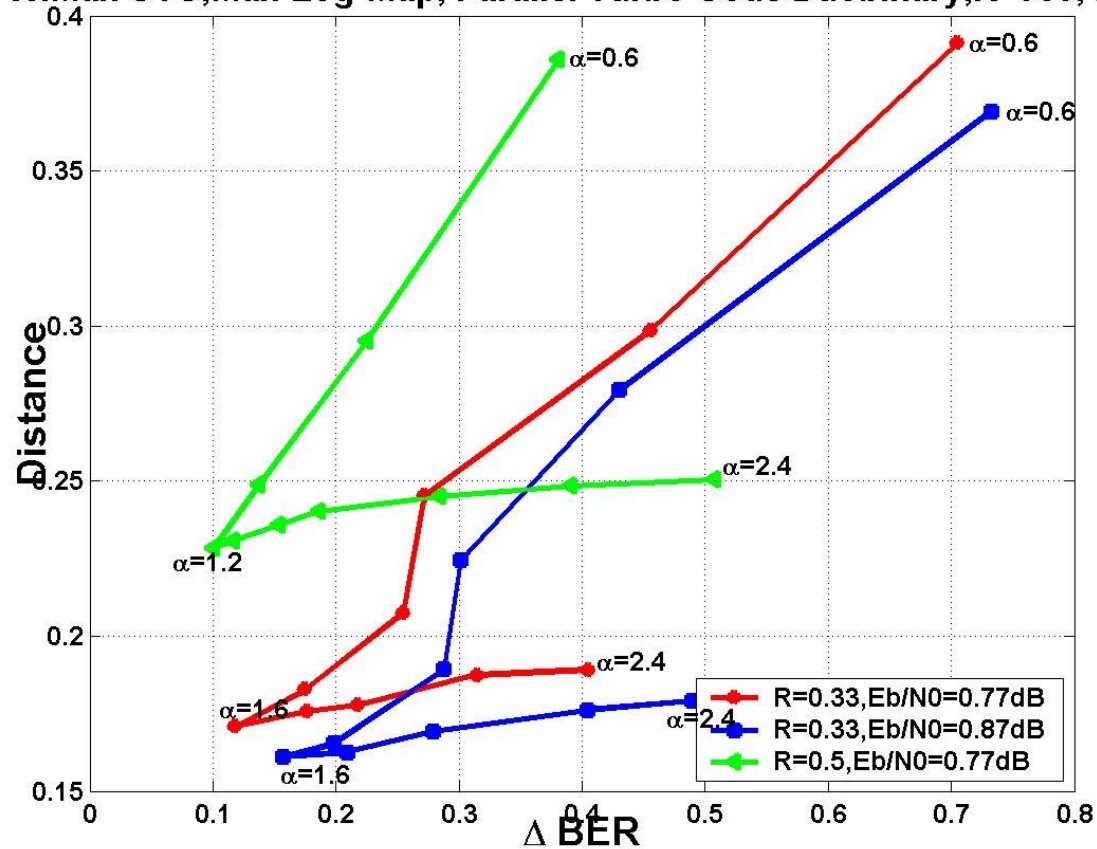


- IF  $\alpha$  too small => ERASURE
- IF  $\alpha$  too big => SATURATION (HARD DECODER)

THERE IS AN OPTIMAL VALUE OF  $\alpha$

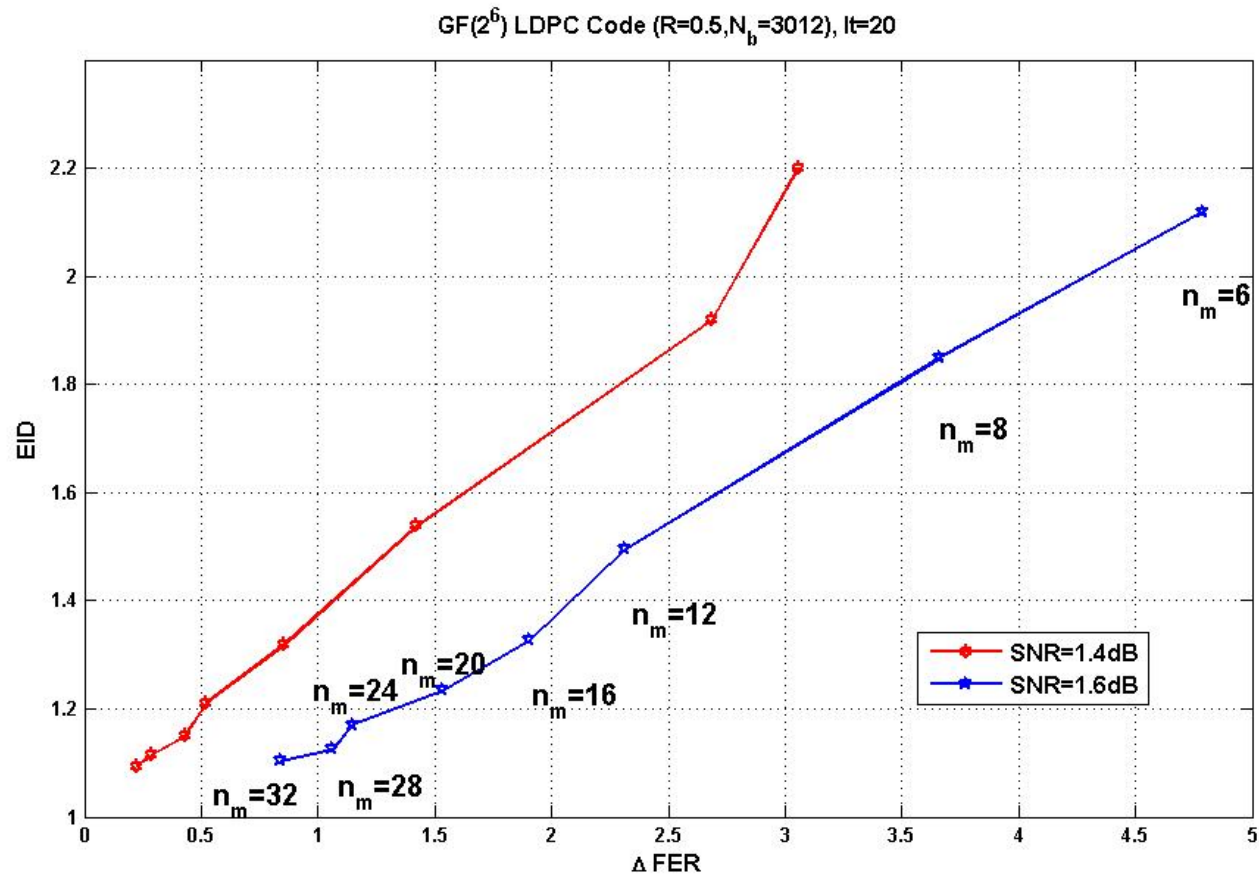
# SIMULATION RESULT : CASE 1

WiMax CTC,Max-Log-Map, Parallel Turbo Code Duobinary,K=960, It=7



# SIMULATION RESULT : CASE 2

- NON BINARY LDPC DECODER ON GF(64) USING EMS ALGORITHM



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# Software simulation

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Three methods to reduce the simulation time:

a) code optimization

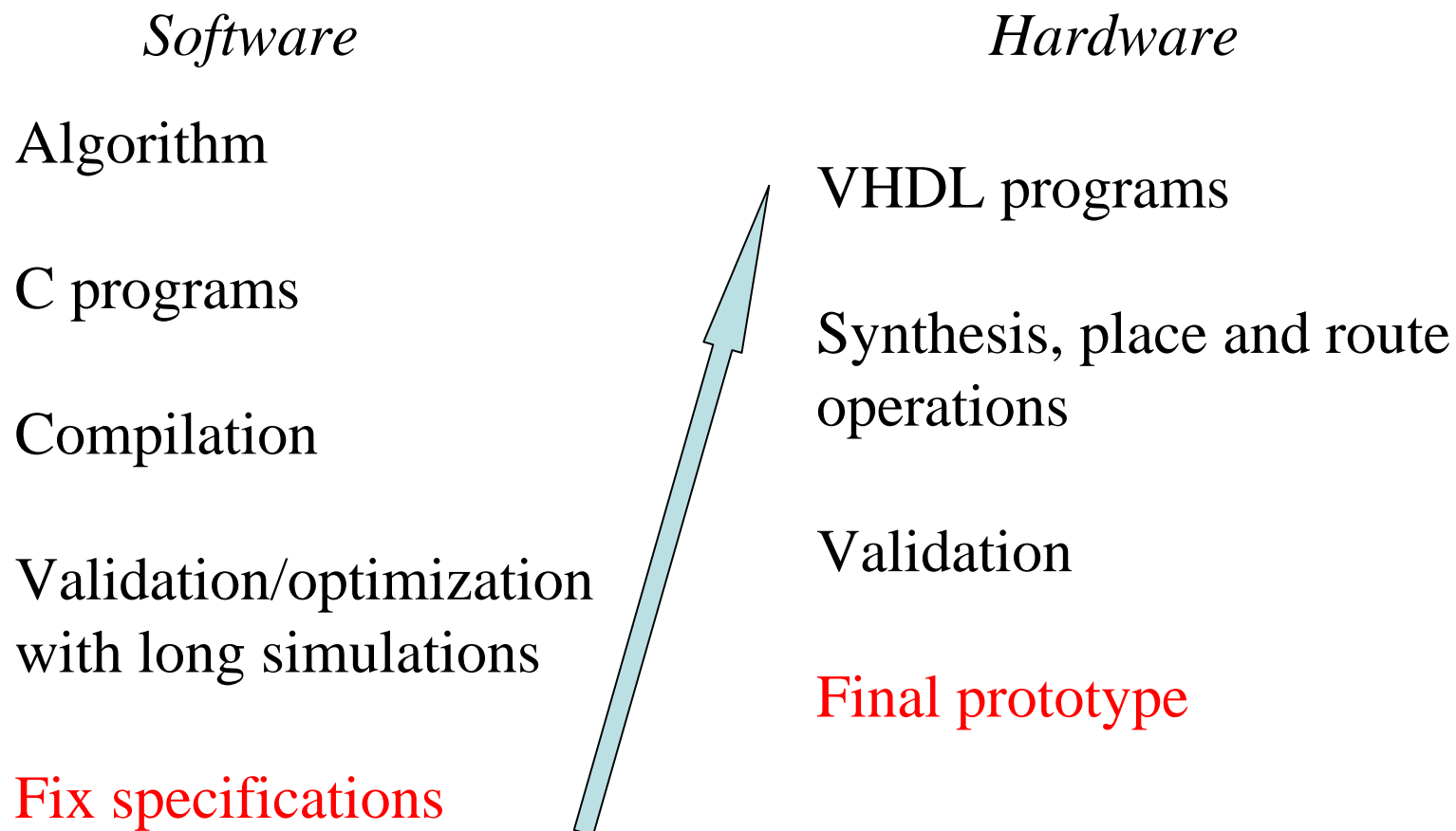
b) powerful computing

c) parallel computing

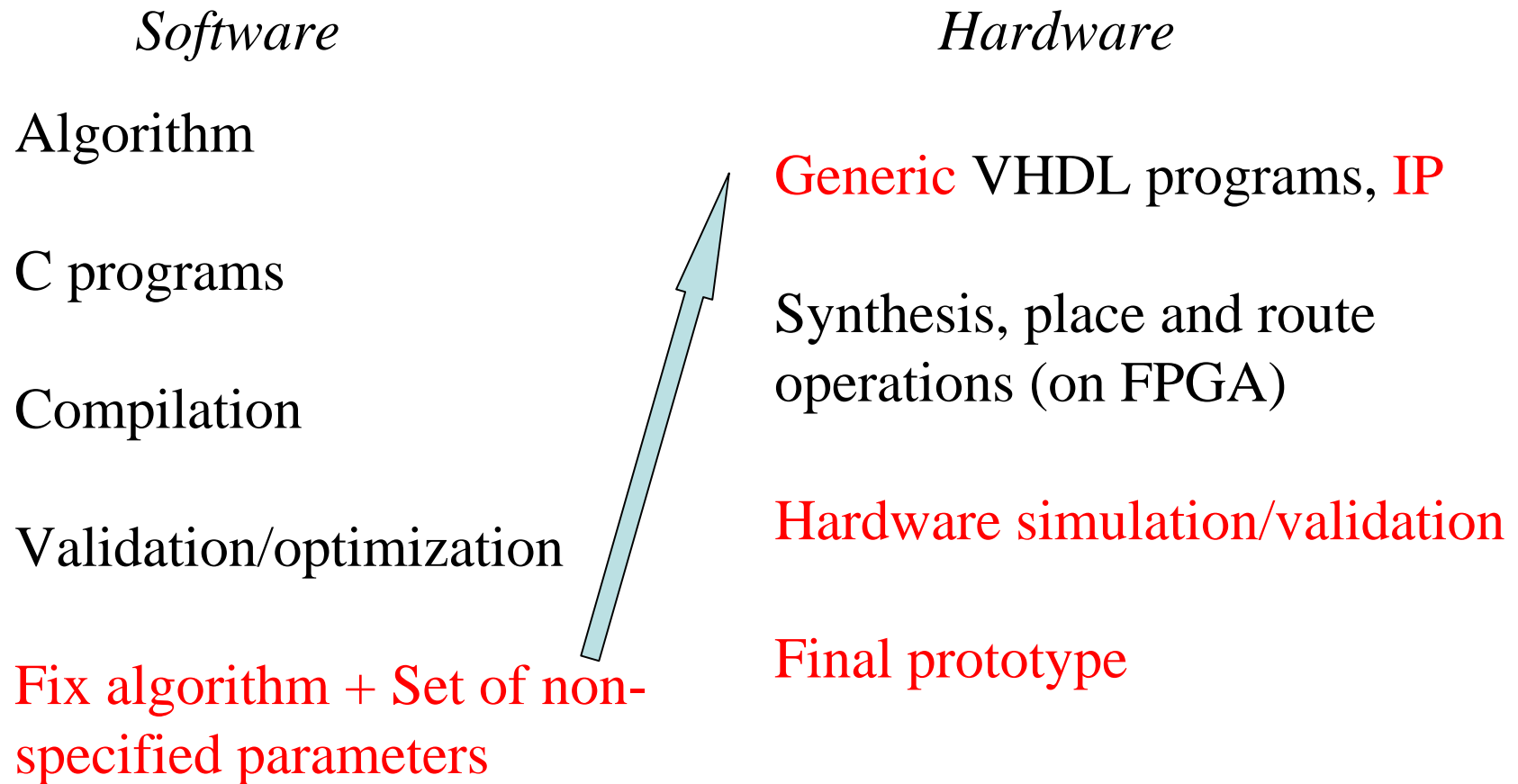
(One Mbps for a turbo-decoder with a cluster of 16 PCs)

**also use hardware emulation**

# Current methodology



# Proposed methodology



# Channel emulation

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Type of communication channel:

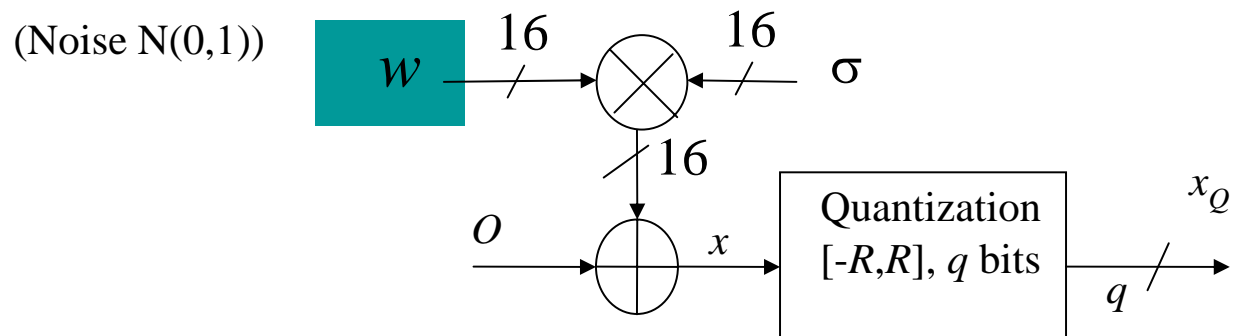
- AWGN
- Rice
- Rayleigh
- ...

All those channels can be derived from Gaussian Noise (with ARMA filter, non-linear operators).

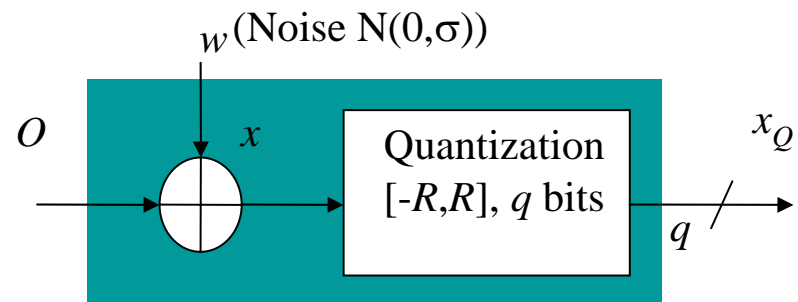
=> Need a White Gaussian Noise Generator (WGNG)

# Nouvelle version

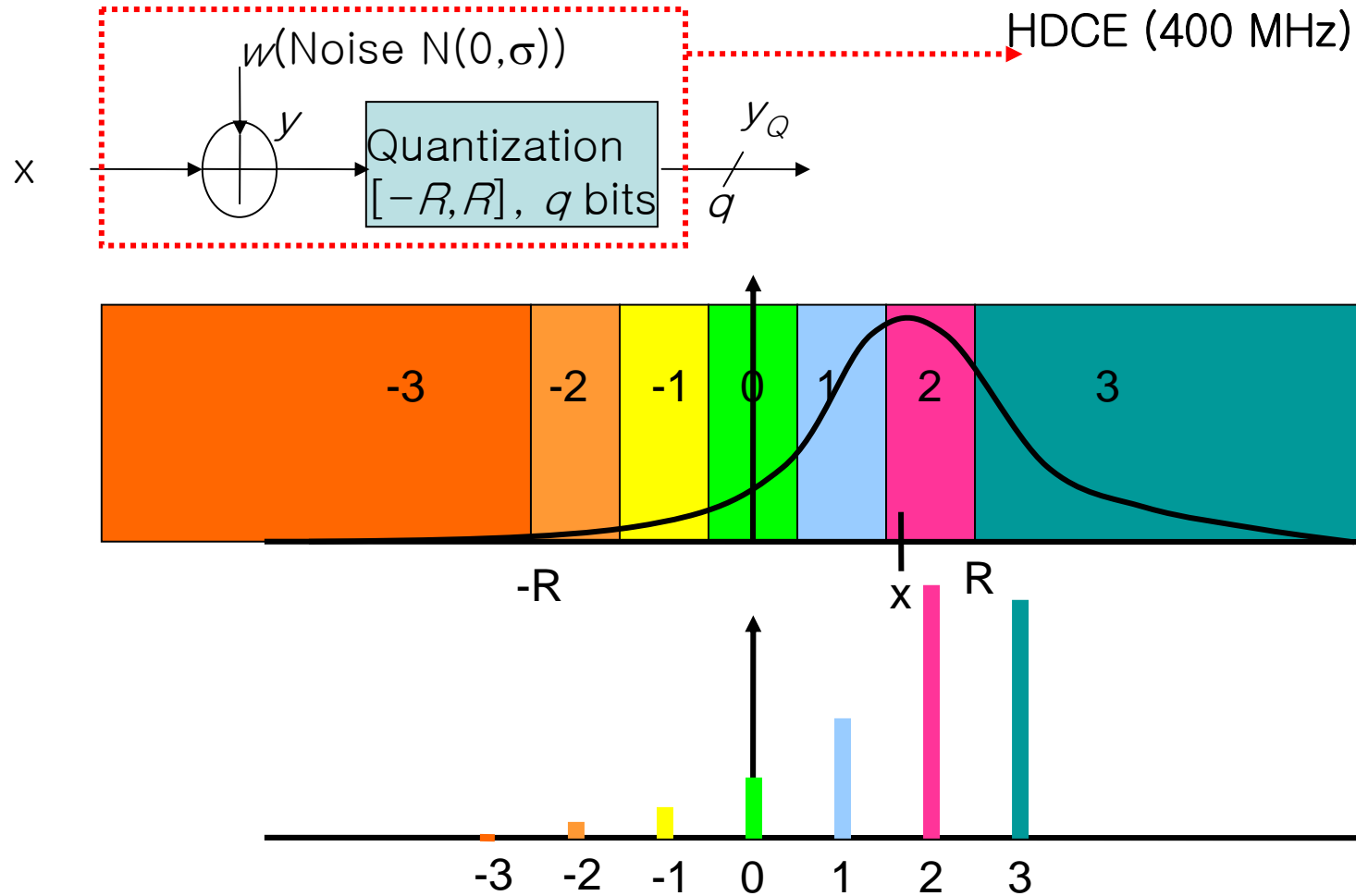
- Old approach



- New approach: emulation of Channel+quantization



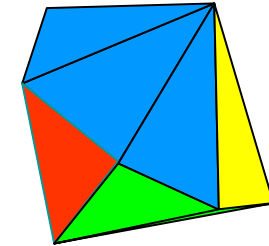
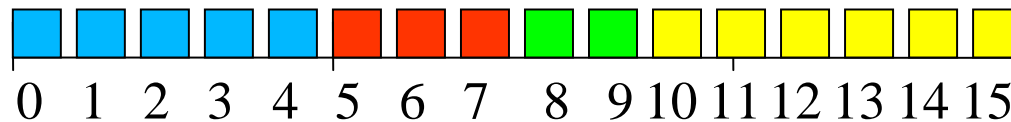
# DISCRETE HARDWARE CHANNEL EMULATOR



- FOR A GIVEN  $x, \sigma, R$ , EQUIVALENT TO A DISCRETE R.V.

# Naive method to generate a discrete variable

- Throw a 16 face dice => it gives a color



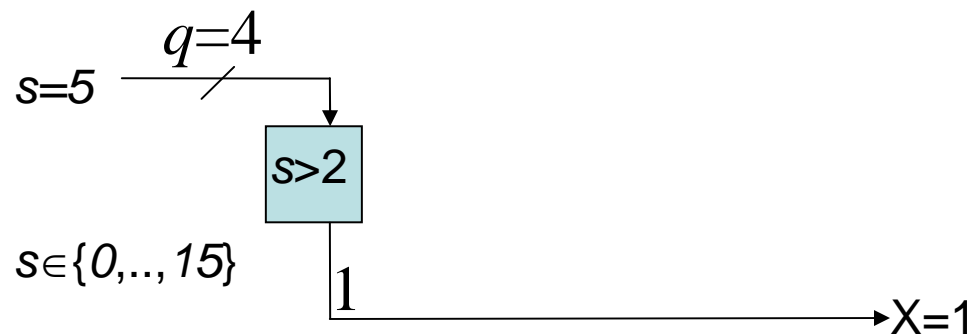
$$\begin{aligned}P(x=B) &= 5/16 \\P(x=R) &= 3/16 \\P(x=G) &= 2/16 \\P(x=W) &= 6/16\end{aligned}$$

**Non uniform random variable !**

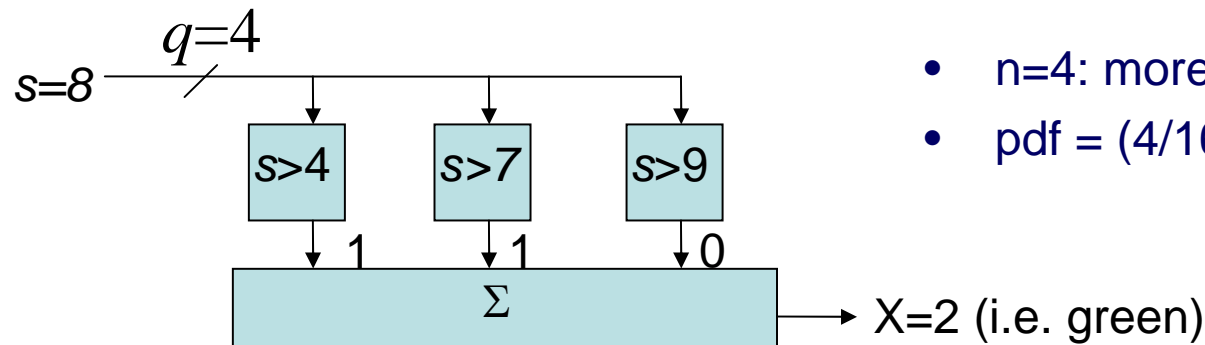
- Associate to each color a number (B=0, R=1, G=2, Y=3)  
=> non uniform distribution.
- By appropriate coloring of the face of the dice, any  $n=4$  value distribution can be obtained (with a precision of  $1/16$ ).

# Naive method to generate a discrete variable

- ELECTRONIC version: generation of a pseudo random number  $s$  between 0 and 15.
- The value of  $s$  is compared to  $n-1$  thresholds to determine the output value.



- $n=2$ : very simple hardware
- Here :  $P(X=0) = 2/16$ ,  $P(X=1)=14/16$



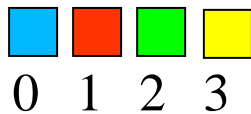
- $n=4$ : more complex hardware
- pdf =  $(4/16, 3/16, 2/16, 7/16)$

- PROBLEM : If  $n=8 \Rightarrow 255$  comparateurs in parallel  $\Rightarrow$  High complexity

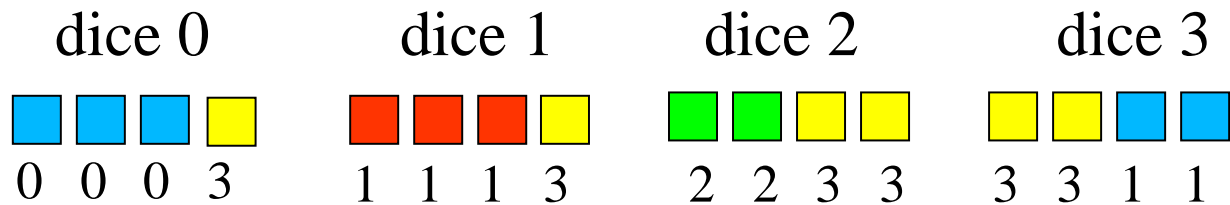
# The Walker 's Alias method

- **Replace a single throw of a 16-face dice with 2 throws of 4-face dices...**

- First throw: Select randomly a color among the 4 (4 different faces).



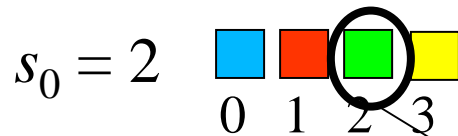
- Second throw: according to the color, select the corresponding two color dice and throw it. Keep the last color (or number).



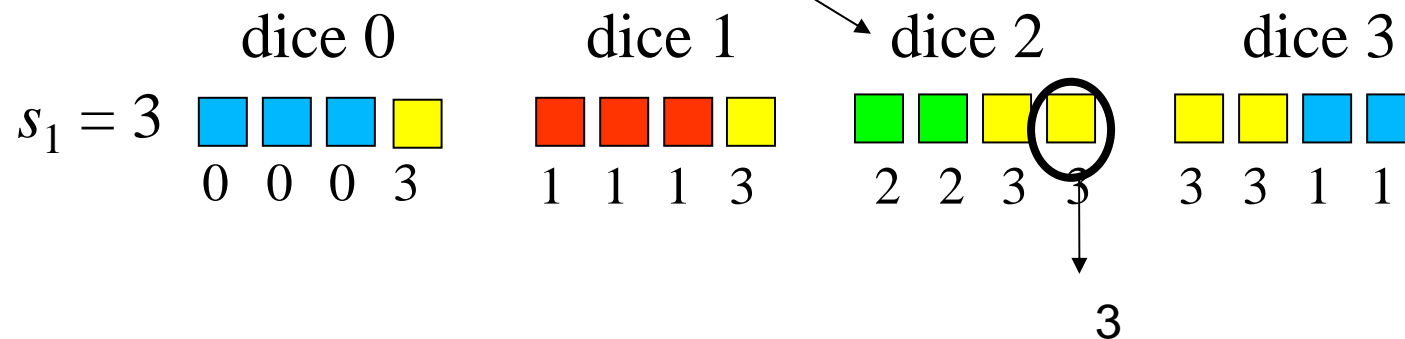
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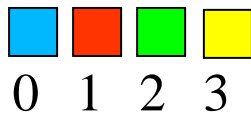
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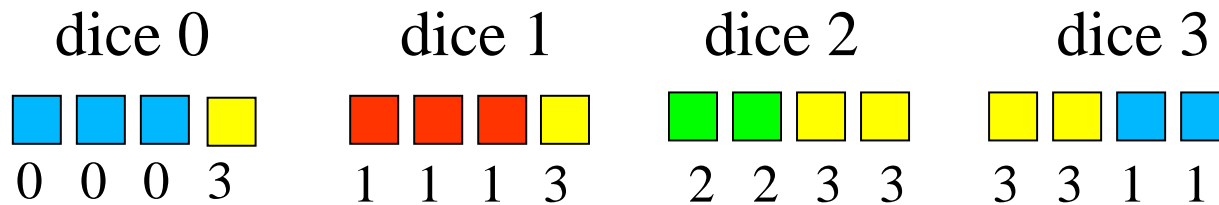
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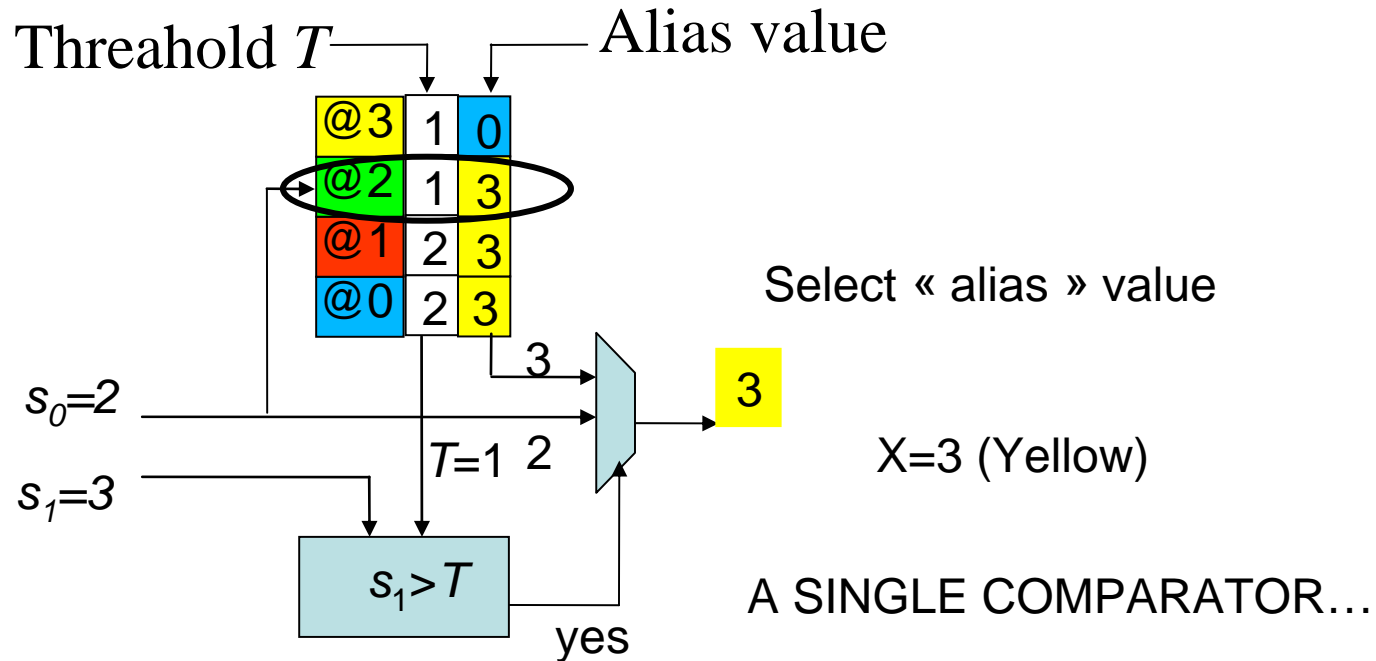
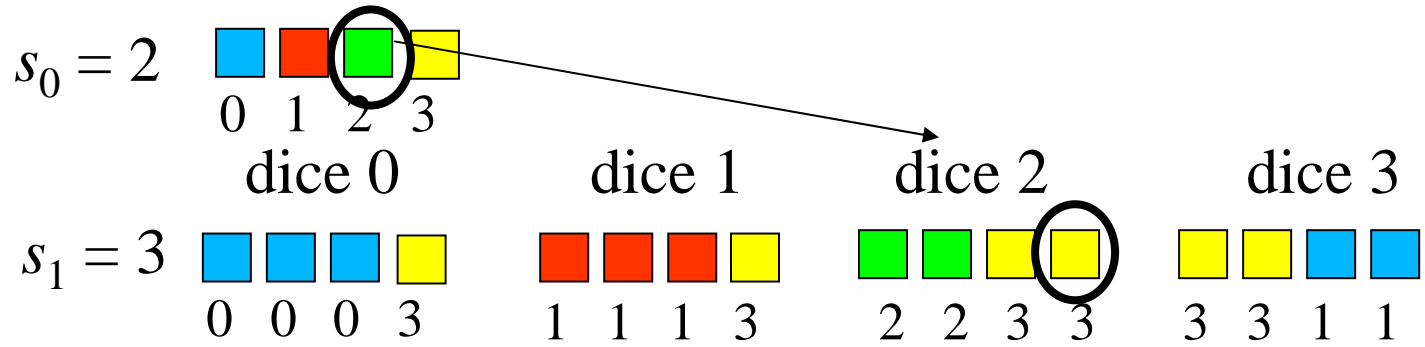


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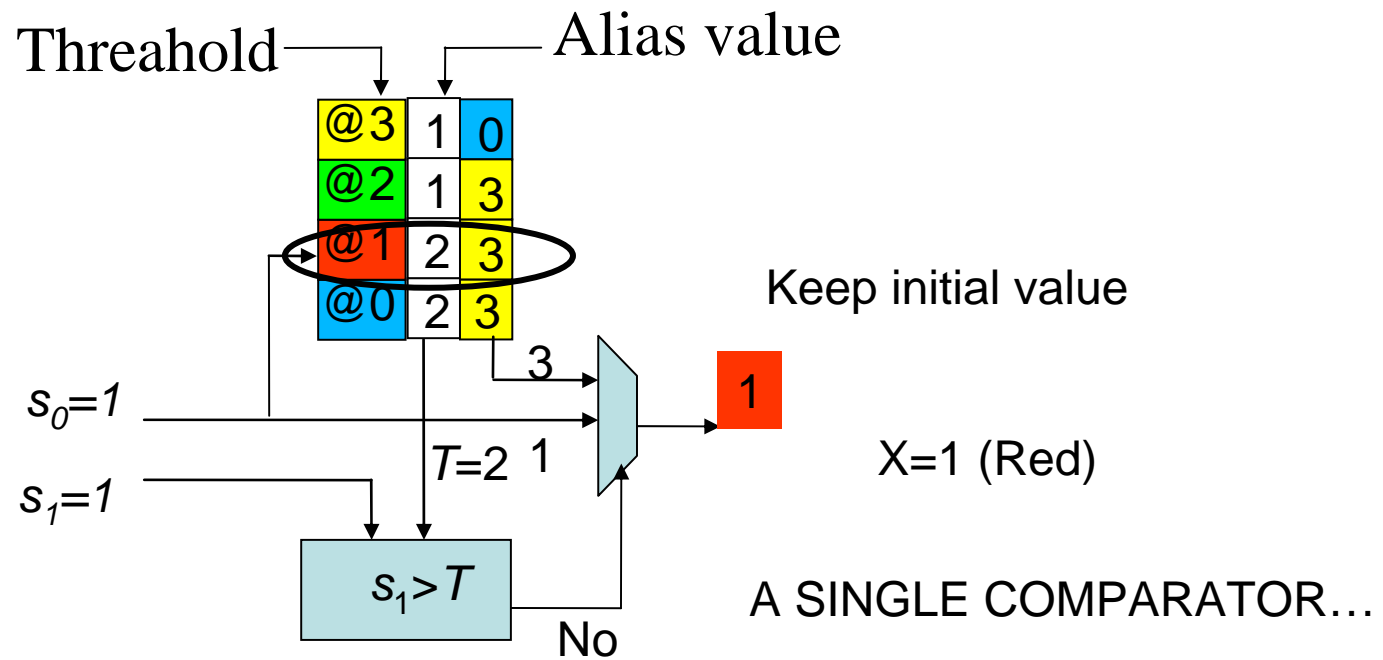
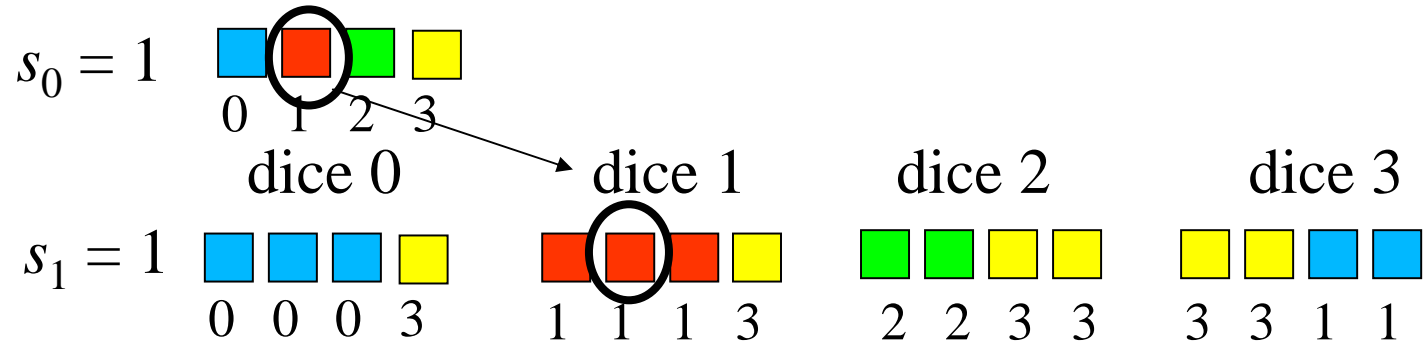
$$P(X=1) = 3/16 \quad P(X=3) = 6/16$$

- We have the initial non-uniform distribution

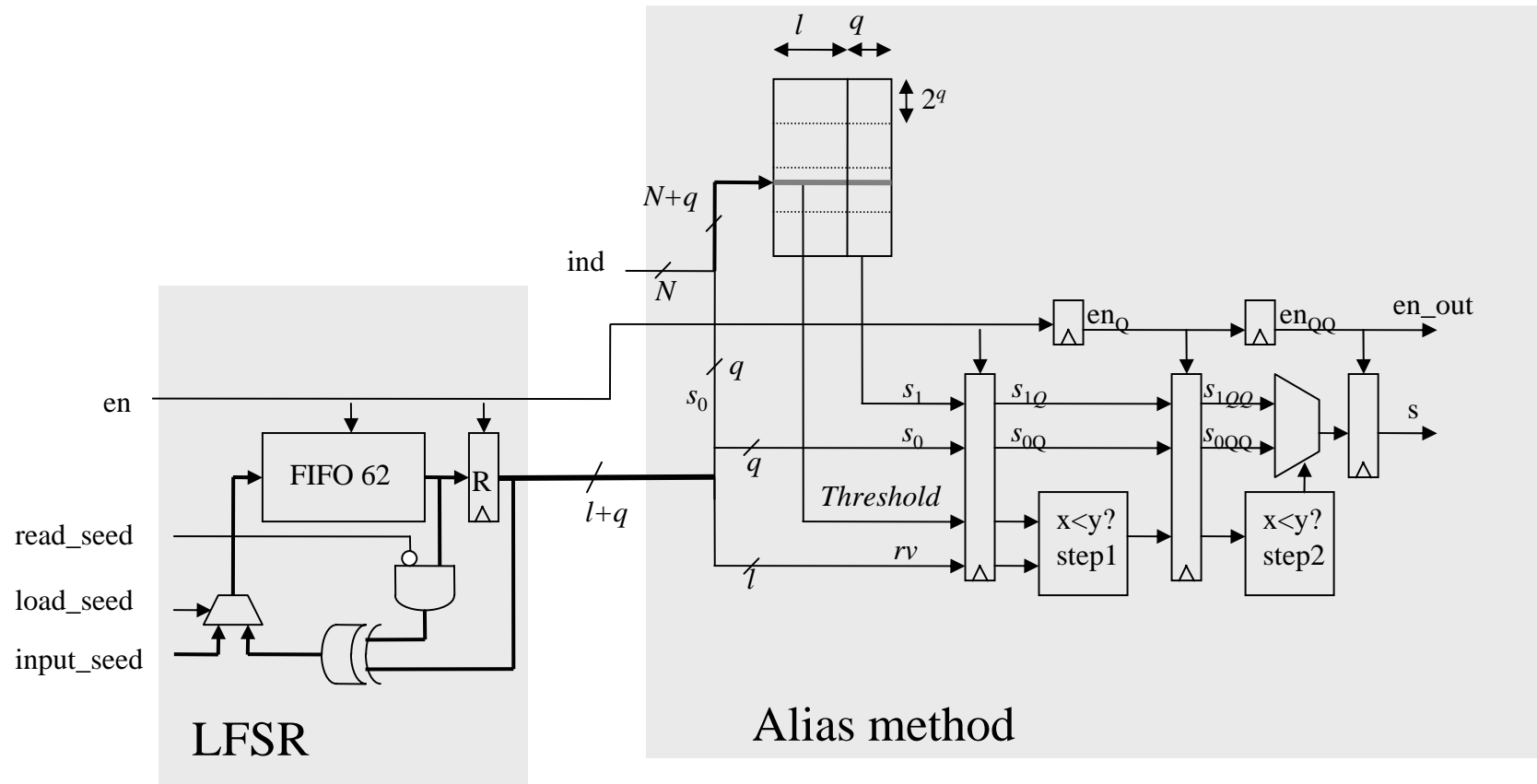
# The Walker 's Alias method



# The Walker 's Alias method



# Architecture of the HDCE



# Functionalities

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- **Generic VHDL**
- **From channel model => optimal generation (in sense of MSE) of alias table.**
  - Ad-Hoc solution
- **Seed management fonctionnality of random generator**
  - Compatible Read/Write process in Hardware and Software
- **C and VHDL fully compatible**
  - Both Observability and high speed simulation.
  - Compatible with the RMCS method

# CONCLUSION

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- **PROPOSITION OF TWO “PRAGMATIC” METHODS TO AVOID “BRUT FORCE” MONTE-CARLO SIMULATIONS**
  - Reduced Monte Carlo Simulation
  - Entropy Inspired Distance
- **THOSE TWO METHODS ALLOW TO TRADE-OFF “TIME OF SIMULATION” VERSUS “PRECISION”**
  - MONTE-CARLO SIMULATION is the “final judge”
  - => HDCE for hardware high speed MC Simulation
- **WE HAVE (WILL) COMBINED THOSE 3 METHODS TO OPTIMIZE THE DECODER ARCHITECTURE OF NON BINARY LDPC CODE OF THE DAVINCI PROJECT.**

# PUBLICATIONS

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- REFERENCES :
- EID :
- A. Singh, A. Al-Ghouwayel<sup>1</sup>, G. Masera, E. Boutillon, "[New Performance Evaluation Metric for Sub-Optimal Iterative Decoders](#)", Accepted to IEEE Communications letters, nov. 2008.
- RMCS :
- E. Boutillon, C. Douillard, G. Montorsi, "[Iterative Decoding of Concatenated Convolutional Codes: Implementation Issues](#)", Transactions of the IEEE, vol. 95, n°6, june 2007.
- <http://www.ict-davinci-codes.eu/>